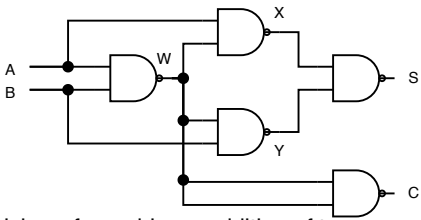


A Level CS C19 Logic Circuits and Boolean algebra

The half adder: a circuit which performs binary addition of two individual bits

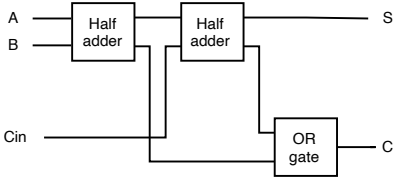
sum bit(S) and **carry bit (C)**

Input		Output	
A	B	S	C
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1



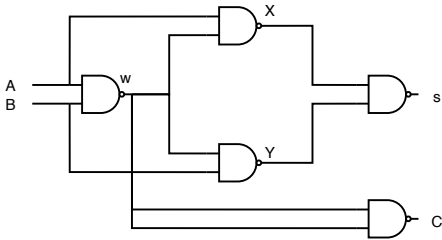
The full adder: a circuit which performs binary addition of two individual bits and an input carry bit. **carry in bit (Cin)**

A	B	Cin	S	Cout
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	0
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1



Boolean algebra: a circuit in which the output is dependent only on the input values

Identity/Law	AND form	IR form
Identity	$1.A = A$	$0 + A = A$
Null	$0.A = 0$	$1 + A = 1$
Idempotent	$A.A = A$	$A + A = A$
inverse	$A.\bar{A} = 0$	$A + \bar{A} = 1$
Commutative	$A.B = B.A$	$A + B = B + A$
Associative	$(A.B).C = A.(B.C)$	$(A+B)+C = A+(B+C)$
Distributive	$A+B.C = (A+B).(A+C)$	$A.(B+C) = A.B + A.C$
Absorption	$A.(A+B) = A$	$A + A.B = A$
De Morgan's	$\overline{A.B} = \bar{A} + \bar{B}$	$\overline{A+B} = \bar{A}.\bar{B}$
Double Component	$\overline{\bar{A}} = A$	$\overline{\bar{A}} = A$



Step 1 NAND Truth Table

NAND Truth Table		
A	B	X
0	0	1
0	1	1
1	0	1
1	1	0

Step 2 Truth Table -> Logic Expression

$$W = \bar{A}.\bar{B} + \bar{A}.B + A.\bar{B}$$

Step 3 Logic Expression X and Y

$$\begin{aligned} X &= \bar{A}.\bar{B} \\ X &= \bar{A}.\bar{B} + \bar{A}.B + A.\bar{B} \\ X &= 0 + 0 + A.\bar{B} \\ X &= A.\bar{B} \\ X &= A + \bar{B} \\ Y &= \bar{A} + B \end{aligned}$$

Step 4 Logic Expression S

$$\begin{aligned} S &= \bar{X}.\bar{Y} \\ S &= \bar{X} + \bar{Y} \\ S &= (A + \bar{B}) + (\bar{A} + B) \\ S &= \bar{A}.B + A.\bar{B} \end{aligned}$$

combinational circuits: a circuit in which the output is dependent only on the input values

Sequential circuit: a circuit in which the output depends on the input values and the previous output

SR flip-flop:

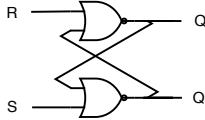
S - set R - reset

set state(self-consistent): $Q = 1, Q' = 0$

unset state: $Q = 0, Q' = 1$

These properties explain why the SR flip-flop can be used as a storage device for 1 bit and therefore could be used as a component in RAM because a value is stored but can be altered.

Input signals		Initial state		Final state	
S	R	Q	Q'	Q	Q'
0	0	1	0	1	0
1	0	1	0	1	0
0	1	1	0	0	1
0	0	0	1	0	1
1	0	0	1	1	0
0	1	0	1	0	1



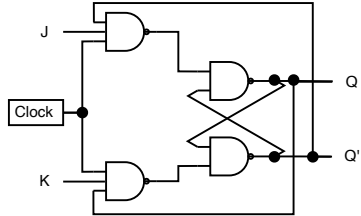
JK flip-flop:

difference: if both J and K are input as a 1 then the values for Q and Q' are toggled (they switch value).

This makes the JK flip-flop a more reliable device because there is no combination of input states that leave uncertainty as to which values are stored.

Input signals Initial state Final state

J	K	Clock	Q
0	0	$\uparrow$	Q unchanged
1	0	$\uparrow$	1
0	1	$\uparrow$	0
1	1	$\uparrow$	Q toggles



difference SR & JK:

1. SR flip-flop has an invalid combination of S and R. SR allow both Q and Q' to have the same value. SR inputs may arrive at different times.
2. JK flip-flop all four combination of values for J and K are valid. JK does not allow Q and Q' to have the same value. JK incorporates a clock pulse for synchronisation

The role of flip-flop in a computer:

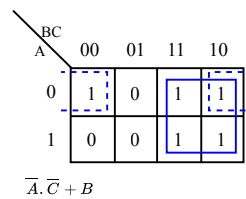
1. A flip-flop can store either a 0 or a 1
2. Computers use bits to store data
3. Flip-flops can therefore be used to store bits (of data)
4. Memory can be created from flip-flops

Karnaugh maps(K-maps): a method of creating a Boolean algebra expression from a truth table.

A K-map can make the process much easier than if you use sum-of-products to create minterms. If applied correctly a K-map produces the simplest possible form for the Boolean algebra expression.

A	B	C	X
0	0	0	1
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

$$\bar{A}.\bar{B}.\bar{C} + \bar{A}.\bar{B}.C + \bar{A}.B.C + A.B.C$$



$$\bar{A}.\bar{C} + B$$

